

# Bounce Rate Optimization

with Markov traffic models

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# Overview

Bounce Rate  
Optimization

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# Background

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## Problem (Optimal tagging)

*In a network of items and tags, how to choose tags for a new item to maximize the probability of a user reaching that item?*

Recommended paper “Optimal tagging with Markov chain optimization” [RG16]

- Model traffic using a Markov chain, and modify transitions in this chain to maximize traffic into a certain state of interest.
- Problem proven to be NP-hard, but with simple greedy approximation from [NWF78]

# Their model

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- Problem formulation:

- 1  $n$  states, plus a new  $\sigma$  state
- 2 2 constrained (details omitted) weight sets:  $q_{ij}$  and  $\bar{q}_{ij}$  - connected and not connected to  $\sigma$  respectively
- 3  $S$  decides which weights to use:  $\rho_{ij}(S) = \begin{cases} q_{ij} & i \in S \\ \bar{q}_{ij} & i \notin S \end{cases}$
- 4 Objective: for a walk  $X$ , find  $S \in [n], |S| \leq k$  maximizing

$$c(S) = \Pr_S[X_t = \sigma \text{ for some } t \geq 0]$$

- Note: only interesting if there are other 'absorbing' states
- Application: system (with ability to estimate network edge weights) recommending tags to user to maximize engagement with their new item

# Theoretical results

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- NP-hard - reduction to vertex cover, but...

## Theorem (Monotonicity)

For all  $S, S' \subset [n]$  with  $S \subseteq S'$ ,  $c(S) \leq c(S')$ .

## Theorem (Submodularity)

(Decreasing marginal value) For all  $S \subset [n]$ ,  $z_1, z_2 \in [n] \setminus S$ ,

$$c(S \cup \{z_1\}) + c(S \cup \{z_2\}) \geq c(S \cup \{z_1, z_2\}) + c(S)$$

- Classic result [NWF78]: monotonic, submodular  
 $\implies (1 - \frac{1}{e})$ -approximation algorithm
- We devise simpler proofs of the two theorems above

# A more realistic approach

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- Our problem formulation:
  - 1  $n$  states plus  $\sigma$ ,  $S$  as before
  - 2 2 differently constrained weight sets:  $q_{ij}$  and  $\bar{q}_{ij}$
  - 3 Notably: user has  $\varepsilon_i$  or  $\bar{\varepsilon}_i$  chance of leaving at each node  $i$
  - 4 Objective: for a walk  $X$ , find  $S$  maximizing  $\ell(S) = \mathbb{E}_S[|X|]$
- Application: given new item with tags, system decides whether or not to feature item in the page for each of the item's tags (decide: connect tag node to new item node?)
- More realistic: system incentivized to maximize *overall* engagement, not just with the new item

# Our model

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- More elusive theoretical guarantees
- In particular, under current conditions: not monotonic (what if new item is *really bad*?)
- We showed with a modified construction it's still NP-hard
- Promising empirical results

# Experimental results

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## ■ Their model vs. our model: expected path length

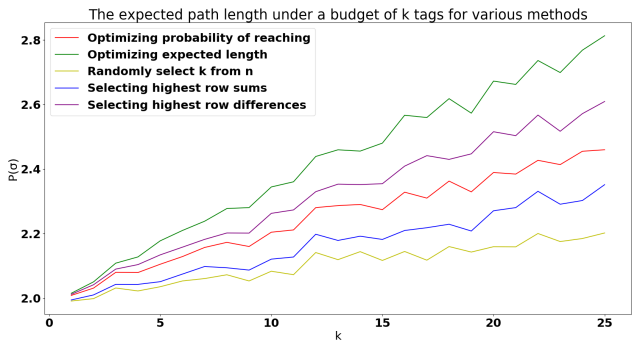


Figure: The expected path length under a budget of  $k$  tags for various methods



# Experimental results

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- Their model vs. our model: probability of reaching  $\sigma$

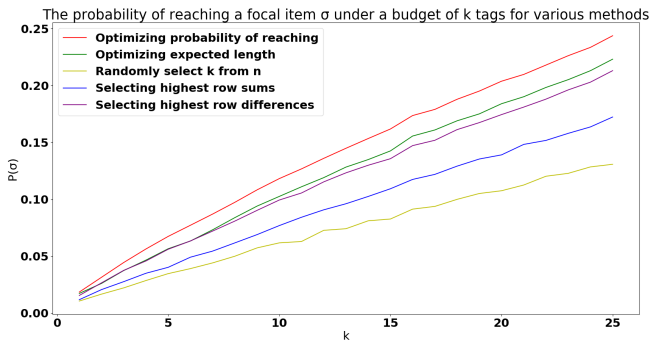


Figure: The probability of reaching a focal item  $\sigma$  under a budget of  $k$  tags for various methods

# Conclusion

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- Even without strong theoretical guarantees yet, *good* empirical evidence
- Model similarity suggests theoretical results around the corner, perhaps for restricted subset of instances
- Additional challenges for incremental computation on large, sparse datasets: LUP decomposition [Rei82]
- Far future: weight estimation techniques, live testing for causal results

# References I

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Ivan Cantador, Peter L Brusilovsky, and Tsvi Kuflik, *Second workshop on information heterogeneity and fusion in recommender systems (hetrec2011)*, ACM, 2011.



George L Nemhauser, Laurence A Wolsey, and Marshall L Fisher, *An analysis of approximations for maximizing submodular set functions—i*, *Mathematical Programming* **14** (1978), no. 1, 265–294.



John Ker Reid, *A sparsity-exploiting variant of the bartels—golub decomposition for linear programming bases*, *Mathematical Programming* **24** (1982), no. 1, 55–69.

# References II

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Nir Rosenfeld and Amir Globerson, *Optimal tagging with markov chain optimization*, Advances in Neural Information Processing Systems, 2016, pp. 1307–1315.

# Appendix: model details

Common setup:

- Markov chain over  $n + 1$  states, initial distribution  $\pi$
- State  $\sigma = n + 1$  is eligible for  $k$  incoming transitions
- Decide  $S \subseteq [n]$  of  $k$  states to “siphon” edge weights to  $\sigma$
- Edge weights specified by  $q_{ij}, \bar{q}_{ij}$  (connected vs. not connected):

$$\rho_{ij}(S) = \begin{cases} q_{ij} & i \in S \\ \bar{q}_{ij} & i \notin S \end{cases}$$

Their formulation:

- Constraint (leeching):  $\forall j \neq \sigma, i, q_{ij} \leq \bar{q}_{ij} \quad \forall i, \bar{q}_{i\sigma} = 0$
- Objective: for  $X_t$  the state of the chain at time  $t$ , find  $S$  maximizing

$$\Pr_S[X_t = \sigma \text{ for some } t \geq 0]$$

# Appendix: model details

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- Markov chain over  $n + 1$  states, initial distribution  $\pi$
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## Our formulation:

- Constraint (leeching):  $\forall j \neq \sigma, i, q_{ij} \leq \bar{q}_{ij} \quad \forall i, \bar{q}_{i\sigma} = 0$
- Constraint (local improvement):  $\forall i, \sum_{j=1}^{n+1} q_{ij} \geq \sum_{j=1}^n \bar{q}_{ij}$
- Constraint (death rate):  $\sum_i^n q_{ij} \leq 1$
- Objective: for  $X_t$  the state of the chain at time  $t$ , find  $S$  maximizing

$$\Pr_S[X_t = \sigma \text{ for some } t \geq 0]$$

# Appendix: datasets

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- Synthetic:
  - $\pi$  selected from normalized  $X \sim \mathcal{U}[0, 1]^n$
  - For  $i \in [n]$ , each  $\bar{q}_i$  also drawn from normalized  $X \sim \mathcal{U}[0, 1]^n$ , then scaled according to  $Y \sim B(6, 3)$
  - Each  $q_i$  is  $\bar{q}_i$  scaled according to another  $Y' \sim B(6, 3)$
  - $q_{i\sigma}$  drawn uniformly so that total of  $q_i$  is between total with  $\bar{q}_i$  and 1.
- Future work: generate data using Last.fm, Delicious, MovieLens datasets from HotRec 2011 workshop [CBK11]